

Identification of Sybil Attack on Social Networks

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Abstract - Social networks play a vital role in daily life since it is vulnerable to many security attacks, namely Sybil attacks. The Sybil attack is an attack where a single user can create a many bogus identity to impersonate like others. The Sybil node in social networking is used for criminal activities such as stealing legitimate information about the user present in social networks it will lead to system degradation process. Since the Sybil identification algorithm does not provide a complete solution to detect the Sybil node in social networks. In order to overcome these drawbacks Sybil defender algorithm gets deployed in a proposed framework to detect the Sybil node in social networks. The Sybil defender algorithm will perform a limited number of random walk on social networks. Sybil defender is a combination of both Sybil identification algorithm and Sybil community detection algorithm. A Sybil identification algorithm is used to detect the Sybil node and Sybil community detection algorithm is used to detect Sybil community around the Sybil node in the social networks. By comparing with the existing approaches Sybil node will be effectively detected using Sybil defender algorithm. This proposed and existing works are measured in terms of evaluation metrics namely non-trustworthy rate, detection rate, packet loss, end to end delivery.

Keywords - Sybil attack, Sybil identification algorithm, Sybil defender algorithm, Social networks, Sybil community detection algorithm.

I. INTRODUCTION

The term social networks refer to the expression of a social relationship, certified or achieved, among individuals, families, households, villages, communities, regions, and so on. Each of them can play dual roles, acting both as a unit or node of a social network as well as a social actor. A Sybil attack is one in which a malicious node on a networks illegitimately claims to be several different nodes simultaneously [1].

The Sybil attack is well known in the context of peer-to-peer, wired, and wireless networks. A Sybil can delay all the messages by a forward lookup to an incorrect or non-existent peer. Finally, it can send

false responses to the receiver. In Sybil attack, an attacker introduces itself in the networks with lots of identities, if an attacker gets large networks identities it can control a large portion of the networks [2].

The number of identities that an attacker can generate depends solely on the attacker's capabilities, which are limited by the bandwidth required for responding to concurrent requests by other peers in the system, the memory required for storing routing information of other peers corresponding to each and every generated Sybil identity, and computation resources required for serving concurrent requests without noticeable delay [3].

To illustrate how this attack works in real systems, imagine a recommender system built over a peer-to-peer overlay in such a system, the goal is to filter information that is likely to be the interest of users based on others' recommendations. In that context, an attacker who can act as multiple users by faking multiple identities can easily out-vote legitimate users' votes on legitimate objects that are subjected to voting. This is almost guaranteed, given that the number of legitimate users who normally vote is always no more than 1% of the total number of users in any realistic recommendation system [4].

In the existing technique, it will detect the Sybil node by using Sybil identification algorithm since this algorithm does not provide complete solution for detecting Sybil node effectively. To overcome the weaknesses of existing technique, in this paper is proposed a Sybil defender which is a centralized Sybil defense mechanism. It consists of a Sybil identification algorithm to identify the Sybil nodes and a Sybil community detection algorithm to detect the Sybil community surrounding a Sybil node. By using these two algorithms, the numbers of attack edges are limited in the social networks [5].

The main contributions of this work includes: Based on performing a limited number of random walks within the social node, this proposed Sybil defender algorithm is more efficient than previous techniques.

The rest of the paper is organized as follows: Section2 describes related works, Section 3

describes the proposed framework, Section 4 describes experimental results and discussion and Section 5 describes the conclusion.

II. RELATED WORKS

In this section, it proposes a survey about how to detect a Sybil attacks present in a social networks and the mechanisms.

Ankush et al., [6] proposed a parental control algorithm to detect a Sybil attack in peer to peer networks based on the reputation scheme since this algorithm applicable only for static networks.

Guojun Wang [5] is applied a Sybil identification algorithm that weeds out all Sybil peers by using neighbor similarity relationship.

Haifeng Yu [7] is introduced a Sybil guard approach. It will detect a Sybil peer present in social networks by generating a random path using pre-computed permutation.

H. Yu, P. Gibbons, M. Kaminsky, and F. Xiao [8] have proposed a Sybil limit. In this method a number of Sybil nodes accepted is reduced by a factor of minus (radicn).

N. Tran et al., [9] have proposed a Gatekeeper is another decentralized Sybil defense scheme that heavily relies on the assumption that the social networks are random expander. This is a strong assumption that has not been validated by previous research. This evaluation shows that Gatekeepers suffers from high false positive and negative rates and cannot effectively identify Sybil nodes on the real-world asymmetric social topologies.

G. Danezi et al., [10] have introduced Sybil Infer, a centralized Sybil defense algorithm, leverages a Bayesian inference approach that assigns a Sybil probability, indicating the degree of certainty, to each node in the networks. It achieves low false negatives at the cost of high computation overhead.

A.Kurve and G. Kesidis [11] proposed Sybil detection via distributed sparse cut in this method. It identifies attack edges and quarantine Sybil clusters. This method works well with dynamic trust graphs as nodes do not need to store any pre-computed data.

C.Hota [12] is proposed a safeguard algorithm, here some arbitrary verifiers are chosen. Each verifier verifies a group of arbitrary nodes, called as suspicious group, by finding paths to each suspicious node and the connection of paths is taken. After connection, the nodes remaining are more likely to be Sybil.

G. Kesidis [13] has proposed a Sybil-proof referral system, which is based on multiplicative reputation chains. Using a multiplicative reputation chain, single step and multi-step referrals can made Sybil proof.

Douceur [14] has proven that without the use of central authority, it is not possible for a system to fully defend against Sybil attack. Hence, in the p2p network, which is fully distributed, Sybil nodes cannot be removed completely from the networks.

Samidha et al., [15] proposed Sybil attack detection on p2p networks based on enhanced Sybil-resilient protocol.

Hengkui Wu et al., [16] proposed a bloom filters to detect a Sybil attack in a distributed system using a historical behavior.

In order to overcome the drawbacks of the above techniques, this proposed framework is used to detect the Sybil attack in social networks.

III. PROPOSED FRAMEWORK

The main objective of the proposed work is to detect Sybil attack which presents in social networks. In the proposed framework, Sybil defender algorithm gets deployed to detect a Sybil attack. Sybil defender algorithm it will limit the number of random walk in a social networks. In Sybil defender algorithm consists of the following two algorithms such as Sybil identification algorithm, a Sybil community detection algorithm, and two supporting approaches which limiting the number of attack edges.

The main task of the Sybil identification algorithm is to determine whether a suspect node is Sybil or not in the social networks. Then, it shows how to efficiently detect the Sybil community around a Sybil node using Sybil community detection algorithm

The purpose of the Sybil community detection algorithm is simply supervising all the nodes in the social graph to find the Sybil community is impractical. Finally, both algorithms are built upon the deduction that the number of attack edges is limited. It is shown in Fig 1.It includes two tasks, namely [17]

1. Detection of attack
2. Sybil defender
 - A.Sybil identification algorithm
 - B.Sybil community detection algorithm

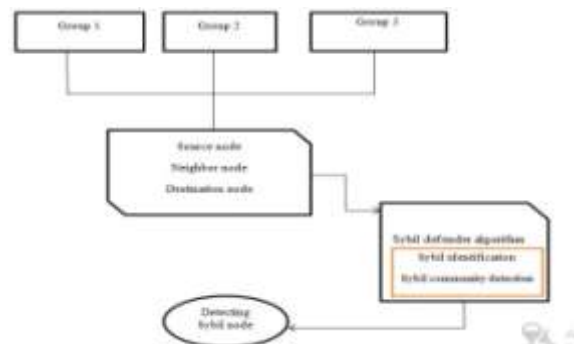


Fig 1. Proposed Framework for detecting Sybil attack in social networks

1. Detection of attack

In this attack is going to identify the type of attacks which is present in the networks. Naturally attacks can be classified into two types:

Active attacks: It will modified all the information transfer between sender and receiver

Passive attacks: It simply listens to all incoming and outgoing messages transfer between sender and receiver, i.e., eaves- dropping, but doesn't harm the system. A peer can be in passive mode and later in active mode [18].

2. Sybil defender

In the Sybil defender algorithm, it is a combination of both Sybil identification algorithm and Sybil community detection algorithm

A. Sybil identification algorithm

A Sybil identification algorithm that takes the social graph $G(V, E)$, a known honest node h , and a suspect node u as input, and outputs whether u is Sybil or not. This algorithm is based on random walks. A sequence of moves of a particle between nodes of G is term as random walk. If the particle is at node i with degree d_i , then the probability that the particle follows the edge (i, j) and moves to a neighbor j is $1/d_i$.

The main idea behind this Sybil identification algorithm is that, as there is a small cut between the honest region and the Sybil region, the random walks originating from a Sybil node tend to get “trapped” into the Sybil region. Also, because it assumes that the size of the Sybil region is not comparable to the size of the honest region.

The number of nodes traversed by the random walks originating from an honest node will be larger than the number of nodes traversed by the random walks originating from a Sybil node, as long as the random walks are long enough to exhibit the difference between the Sybil region and the honest region, and it performs the random walks many times. For simplicity, it defines the number of times one node being traversed by a set of random walks as the frequency of that node [18].

Fig 2 shows a Sybil identification algorithm for identifying the Sybil node in the social networks. Here step 1 to 11 will performs a preprocessing process step 12 to 26 will perform random walk on social node. After completing all process Sybil node get detected.

```

1. j={h}
2. For i=1 to f do
3. Perform a random walk with length ls=logn originating from h
4. J=U (the ending node of the random walk)
5. End for
6. L=l_min
7. While l<=l_max do
8. For i=j.first() to J.last() do
9. Perform R random walks with length l originating from i
10 Get ni as the no of nodes with frequency no smaller than t
11. End for
12. Output(l,mean({ni i∈J}),stdDeviation({ni i∈J}))
13. L=l+100
14. End while
15. L=l_0
16. While l<=l_max do
17. Perform R random walks with length l originating from k
18. M=the number of nodes whose frequency is no smaller than t
19. Let the tuple corresponding to length l in the outputs of algorithm 1 be (l, mean, stdDeviation)
20. If mean - m > stdDeviation * α then
21. Output k is Sybil
22. End the algorithm
23. End if
24. L=l*2
25. End while
26. Output k is honest
    
```

Fig 2 shows a Sybil identification algorithm

B. Sybil community detection algorithm

After one Sybil node is identified, The Sybil community detection algorithm can be used to detect the Sybil community surrounding it. The Sybil community detection algorithm takes the social graph $G(V, E)$ and a known Sybil node as input, and outputs the Sybil community around us. The Sybil nodes can be identified by using Sybil identification algorithm or any previous scheme. It defines a Sybil community as a subgroup of G consisting of only Sybil nodes, and there is no small cut in this sub graph.

The reason it makes this definition is that if a small cut does divide the Sybil region into two parts $S1$ and $S2$, and the known Sybil nodes is s in $S1$, then, from the point of view of us, the honest region and $S2$ are similar, because there is already a small cut between $S1$ and the honest region and also a small cut between $S1$ and $S2$. When there is a small cut in the Sybil region, this algorithm can detect the Sybil community s .

This algorithm based on performing partial random walks originating from s . Each partial random walk behaves the same as the simple random Walks used in the Sybil identification algorithm, except that it does not traverse the same node more than once.

Therefore, when a partial random walk reaches a node with all the neighbors traversed by itself, this partial random walk is “dead” and cannot proceed.

This property makes a partial random walk originating from a Sybil node less likely to leave the Sybil region, compared with a simple random walk, because many such walks “die” when they hit the border of the Sybil region. Similar to the Sybil identification algorithm, the intuition behind this algorithm is that the partial random walks originating from a Sybil node tend to be trapped within the Sybil region, and thus, it can detect the Sybil community by examining the nodes traversed by the partial random walks. Fig 3 shows a Sybil community detection algorithm for identifying the Sybil community around a Sybil node in the social networks.

In Sybil community detection algorithm will be used to perform a walk length between node here step 1 to step 8 perform random walk and step 9 to step 13 output of walk length estimation[18].

```

1. i=0
2. Deadwalknum=0
3. While deadwalknum<0 do
4. i=i+1
5. Deadwalknum=0
6. For i=1 to do
7. Perform a partial random walk originating from s with length l if the partial random walk is dead before it reaches
8. Deadwalknum++
9. End if
10. End for
11. Deadwalknum=deadwalknum
    
```

Fig 3 shows a Sybil community detection algorithm

Finally, Sybil attack is detected on social networks by using above two algorithms namely Sybil identification algorithm and Sybil community detection algorithm.

IV. EXPERIMENTAL RESULTS AND DISCUSSION

Performance Metrics

The evaluation metrics are mainly involved to calculate the effectiveness of the performance; the performance of the proposed framework is measured in terms of the quality measures namely

1. Non-trustworthy rate
2. Detection rate
3. Packet Delivery Ratio (PDR)
4. End-To-End Delay

Non-trustworthy rate

Non-trustworthy rate is the ratio of the number of honest peers which are erroneously marked as a Sybil peer to the number of total honest peers.

$$\text{Non-trustworthy} = \frac{\text{No.of honest peer marked as Sybil}}{\text{No.of total honest peer}}$$

Detection rate

Detection rate is the proposition of detected Sybil/malicious peers to the total Sybil/malicious peers.

$$\text{Detection Rate} = \frac{\text{Detected Sybil peer}}{\text{Total Sybil/malicious peer}}$$

Packet Delivery Ratio (PDR)

The ratio of the number of data packets receives the packet data send to the destination. This illustrates the level of delivered data to the destination.

$$\text{PDR} = \frac{\sum \text{Number of packets receive}}{\sum \text{Number of packets sends}}$$

End-To-End Delay

The average time it takes a data packet to reach the destination. This includes all possible delays caused by buffering during route discovery latency, queuing at the interface queue. This metric is calculated by subtracting time at which first packet was transmitted from source from time at which first data packet arrived to destination.

$$\text{End to end delay} = \frac{\sum (\text{arrive time-send time})}{\sum \text{Number of connection}}$$

Existing technique:

In existing framework, Sybil node gets detected by using Sybil identification algorithm based on neighbor similarity relationships. In this approach group of all the nodes which has similar behavior using similarity trust relationship hence it will act like an identifier source. They can send Identifiers to others as the system regulates. If a node sends less or more, the system can be having a Sybil attack node. The performance of existing framework is measured using the above metrics [18].

The below table shows a non-trustworthy rate detection and here number of inputs will be a node which is deployed. The table also shows that when number of honest peer get marked as Sybil increases than non-trustworthy also get increases.

The results were shown in the Table I. This table shows the performance of the non-trustworthy rate detection, which is detected using Sybil defender algorithm and Sybil identification algorithm.

TABLE I. Non-trustworthy rate using different inputs

No of nodes	Technique used			
	Existing technique		Proposed technique	
	Sybil identification algorithm		Sybil defender algorithm	
	No of honest peer	Non-trustworthy rate	No of honest peer	Non-trustworthy rate
20	12	0.56	14	0.4
30	24	0.69	26	0.57
40	31	0.79	34	0.68
50	38	0.68	42	0.5
60	52	0.9	54	0.87

The below graph Fig 4 shows a non-trustworthy rate detection, which is detected using Sybil defender and Sybil identification algorithm with a number of inputs. Here x-axis denotes number of nodes deployed and y-axis denotes non-trustworthy rate detection. By observing this graph Sybil defender produces low non-trustworthy rate detection.

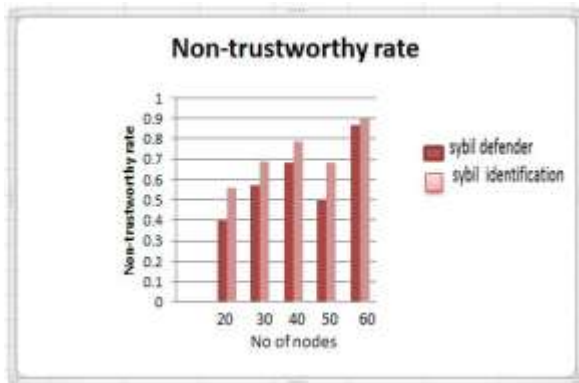


Fig 4. Graph for non-trustworthy rate using different inputs

The below table shows the performance of the detection rate, which is detected using Sybil defender algorithm and Sybil identification algorithm. The results were shown in the Table II.

TABLE II. Detection rate using different inputs

No of nodes	Technique used					
	Existing technique			Proposed technique		
	Sybil identification algorithm			Sybil defender algorithm		
	Total sybil peer	Undetected sybil peer	Detection rate	Total sybil peer	Undetected sybil peer	Detection rate
20	12	4	0.5	10	9	0.75
30	15	8	0.6	15	13	0.6
40	26	15	0.63	30	10	0.63
50	35	21	0.789	29	20	0.789
60	46	28	0.85	38	28	0.85

The below graph Fig 5 shows a detection rate, which is detected using Sybil defender and Sybil identification algorithm with a number of inputs. Here x-axis denotes number of nodes deployed and y-axis denotes detection rate. By observing this graph Sybil defender produces high detection rate

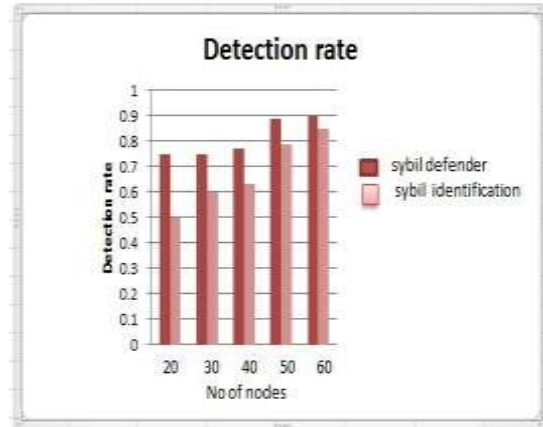


Fig 5. Graph for detection rate using different inputs

The below table shows the performance of the packet loss, which is detected using Sybil defender algorithm and Sybil identification algorithm. The results were shown in the Table III.

TABLE III. Packet loss using different inputs

No of nodes	Technique used							
	Existing technique				Proposed technique			
	Sybil identification algorithm				Sybil defender algorithm			
	No of packet sends	No of packet receive	Packet delivery	Packet loss	No of packet sends	No of packet receive	Packet delivery	Packet loss
20	20	15	0.75	10	22	11	0.7	7
30	30	13	0.6	20	35	30	0.69	14
40	40	40	0.625	9	37	30	0.64	5
50	50	50	0.61	14	55	50	0.70	12
60	60	60	0.5	10	57	55	0.89	8

The below graph Fig 6 shows packet loss, which is detected using Sybil defender and Sybil identification algorithm with a number of inputs. Here x-axis denotes number of nodes deployed and y-axis denotes packet loss. By observing this graph Sybil defender produce low packet loss.

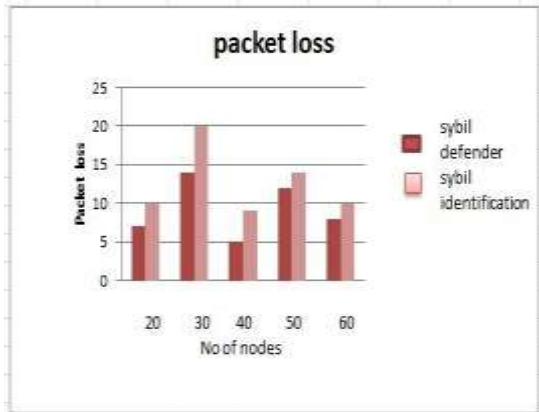


Fig 6. Graph for packet loss using different inputs

The below table shows the performance of the end to end delay, which is detected using Sybil defender algorithm and Sybil identification algorithm. The results were shown in the Table IV.

TABLE IV. End to end delay using different inputs

No of nodes	Technique used					
	Existing technique			Proposed technique		
	Sybil identification algorithm			Sybil defender algorithm		
	Send time	Arrive time	End to End delay	Send time	Arrive time	End to End delay
20	10	15	0.89	14	14	0.7
30	2	4	0.8	9	9.1	0.78
40	1	3	0.79	4	4.2	0.67
50	2	16	0.6	7	7	0.58
60	5	6	0.78	8	8	0.64

The below graph Fig 7 shows end to end delay, which is detected using Sybil defender and Sybil identification algorithm with a number of inputs. Here x-axis denotes number of nodes deployed and y-axis denotes end to end delay. By observing this graph Sybil defender it produces a low end to end delay rate.

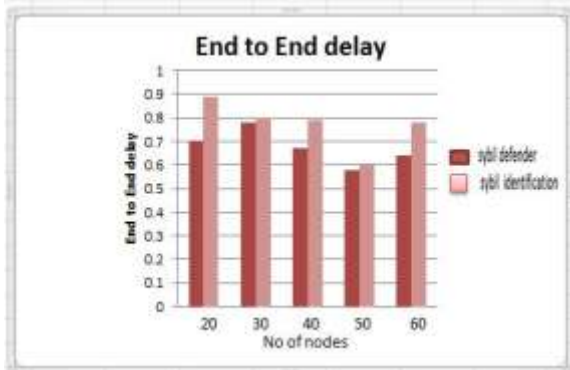


Fig 7. Graph for end to end delay using different inputs

V CONCLUSION

Social networks are vulnerable to security attack, namely Sybil attack. Sybil attacker can create fake accounts in social networks for stealing legitimate user information it can be detected by using Sybil defender algorithm. The Sybil defender algorithm is evaluated in terms of performance metrics namely non-trustworthy rate, detection rate, packet loss and end to end delivery. By comparing with existing technique Sybil defender will be provided an effective result.

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