Clustering Based Dynamic Load Balancing Algorithm

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Abstract: The growth of Internet requires speedy web servers so that reaction time can be reduced and better services can be given to clients. The multiple servers are interconnected to reduce the workload and partitioning the request. There is a vital need of a load balancing technique to distribute this load among web servers by achieving performance and minimizing response time. In this paper, we proposed a framework for load balancing in heterogeneous web server clusters. Load is distributed on the basis of memory and processing requirements. Framework is explained with pseudo code and experimental results are also shown.

Keywords: Web Server Clusters, Load Balancing, Response time, Throughput

I. INTRODUCTION

In a distributed system, jobs arrive at various hosts independently, causing uneven loads on the hosts. To balance the load effectively, there are various factors that need to be considered. Load balancing means to transfer jobs among different processors in a distributed system. Its aim is to keep the performance of application by keeping processoridle time and interprocessor communication as low as possible. To balance load effectively, there are various factors that need to be considered like performance of system after doing load balancing, the algorithm to be used, load index value and impact of load balancing on system behavior [9].

In case, jobs are coming with constant workloads, static load balancing can be used. In such case, there may be one central processor that can act as a pre-processorto the computations. For other applications with inconsistent workloads dynamic load balancers can be used. Several strategies for static and dynamicload balancing have been developed.

Load balancing in distributed systems is a challenge due to the independence of the processors and the IPC overhead incurred. A parallel and distributed systemwhere large task is divided among multiple hosts is the finest option for computing parallel applications. In such systems, there are chances that one node becomes overloaded and others under loaded. The algorithms which help to transfer jobs from heavily loaded nodes to lightly loaded nodes are known as load balancing algorithms.

II. FORMAL DESCRIPTION OF THE PROPOSED ALGORITHM

There is a need for various load balancing algorithms to improve the performance of nodes in distributed systems. In this research paper, we propose a new load balancing algorithm to distribute the load among nodes so that all nodes will be well balanced.

In this algorithm, we consider the following things:

- I) Jobs allocated on the basis of memory requirement and processing speed.
- II) Clusters are defined on the basis of memory partitions.

Following are the various terms defined:

A) About Job Requirements:

J_i: Incoming Job ID

J_{im}: Incoming Job Memory Requirement

 J_{ip} : Incoming Job Processing Speed Requirement

Jis: Job Status

where i = 0 to n (n is total number of jobs)

B) About Servers (Total Number of Servers Say k=6)

 $\begin{array}{l} S_k: \mbox{ Server ID} \\ S_{ck}: \mbox{ Server ID in cluster `C'} \\ S_{ckm}: \mbox{ Server having maximum memory} \\ S_{ckp}: \mbox{ Server having maximum processing} \\ speed \\ S_{ckl}: \mbox{ Memory left after job being allocated} \end{array}$

S_{ckj}: Job Assigned to server

C) About Clusters (Total number of clusters say L=2)

C_L: Cluster ID

C _{Lrs} : Intial va	alue of Range
C _{Lrl} : Last val	ue of Range

Initial Information:

- A) Total number of incoming Jobs0 to n (n=7 and 0 means no job in the queue)
- B) Total number of servers available is k = 6
- C) Defining Range of Clusters (L =2) on the basis of memory partitions are as follows:

Cluster 1	Cluster 2	
$C_{1RS} = 0$	C _{2RS} =1001	
$C_{1RL} = 1000$	$C_{2RL} = 2000$	
D) For current server:		
j=1 to count		
Count is same as k		

Next step is to allocate the jobs on the basis of memory and processing speed requirements. Firstly, job's memory requirement is compared with cluster's range and then job is assigned to particular server under that cluster in round robin manner.

Step 1: Assigning Jobs to servers in the clusters:

Initially, following is the status:

 $\begin{array}{l} J_i = 1 \mbox{ to } n \mbox{ (where } n = 7) \\ T_i = 1 \mbox{ to count (where count = 6)} \\ S_k = 1 \mbox{ to } k \mbox{ (where } k = 6) \\ C_L = 1 \mbox{ to } l \mbox{ (where } L = 2) \\ J_{is} = "Unallocated" \\ \textit{Now for allocating jobs, first Cluster is selected:} \end{array}$

Condition 1: (Selecting Cluster)

if (J_{im} > C_{LRS}) AND (J_{im} < = C_{LRL}) Then Select Cluster C_L Else C_L =+1 (move to next cluster) End if *Now select Server under selected Cluster.*

Condition 2: (Selecting Web Server)

iii) Find status of the server : if ($S_{ckl} < 0$) then

J=+1 AND k=+1 AND check with condition i) and ii) Until satisfied

Finally, when current job is allocated to any server, i = +1 (move to next job). If current job do not get the server, it is queued and checked once again later on.

Step 2: Load Balancing (LB)

Once jobs are allocated, some web servers might be unbalanced. Thus next step is to balance the load among web servers. For doing LB, following things will be considered:

- i) Status of the server i.e. uneven
- ii) Servers with no load

Basic Logic:

Step 1: Check server's status as 'Uneven'

Step 2: To make status even, LB is done

- a) Within the cluster
- b) Between Clusters

Step 3: If at any point of time, status cannot be changed from UNEVEN to EVEN, then

- a) Job would be waited in Queue
- b) When the respected server would be free, reintroduce the job.

Initially,

- $\begin{array}{ll} a) & S_k=1 \mbox{ to } k & // \mbox{ Server ID, } k=6 \\ b) & i=0 \mbox{ to } n & // \mbox{ Jobs, } n=7 \\ c) & C_L=1 \mbox{ to } L & // \mbox{ } L=2 \\ d) & J=1 \mbox{ to count } & // \mbox{ count is jobs allocated to } \end{array}$
- ALGORITHM:

server

Step 1: a) $J_i = 0$ to n // n = 7 for jobs J = 1 $S_k = 1$ to k $C_L = 1$ to L b)IfS_{cks} = 'UNEVEN' Then

 $J=S_{\text{ckj}}\!/\!/$ j contains all the allocated jobs of servers

Step 2: a) Calculate left over memory after removing jobs for 'j'

$$\begin{array}{l} J_i = j \\ S_{ckl} = S_{ckm} + J_{im} \\ b) \quad If (S_{ckl} > 0) \\ Then \\ S_{cks} = "EVEN" \\ J_{is} = "UNALLOCATED" \\ Goto step 3(a) \\ k = +1 \\ goto step 1(b) \\ else if (S_{ckL} < 0) \\ then \\ j = +1 \\ repeat step 2(a) \\ else if (S_{ckL} == 0) \\ k = +1 \\ goto step 1(b) \end{array}$$

Step 3: a) k = +1

// checking for other servers

b) if (k <=max.server in a cluster) then if ($[J_{in} \leq S_{ckm}]$ AND $[J_{in} \leq S_{ckp}]$) then goto step 3(c)else goto step 3 (a) else $C_l = +1$ // move to next clusters K = 1endif Goto step 3(b) c) If ($S_{cks} == "UNEVEN"$) Then Goto step 3(a) Else $S_{ckl} = S_{ckm} - J_{im}$ If ($S_{ckl} < 0$) Then $S_{ckl} = S_{ckm} + J_{im}$ Goto step 3(a) Else $S_{cki} = + J_i$ // Adding jobs to allocated list of servers J_{is} = "ALLOCATED"

III. EXPERIMENTAL RESULTS AND ANALYSIS

The results of the experiments are shown below. Figure 1 presents how the average execution time varied with the total load of the system. Thus it is shown that with the increase in total number of jobs, execution time is also increased at consistent rate. Even if we have more number of jobs, then also execution time rate would be same.

TABLE 1 execution time w.r.t total number of jobs		
Total Number of Jobs	Execution Time (in Secs)	
3	3.73	
5	7.36	
7	9.72	
10	11.09	





IV. CONCLUSION AND FUTURE WORK

We proposed a framework for load balancing in heterogeneous web server clusters. Load is distributed on the basis of memory and processing requirements. Preliminary evaluation reveals that use of this algorithm is necessary to improve the performance of web servers by proper resource utilization and reducing the mean response time by distributing the workload evenly among the servers in the cluster.

In future, the proposed algorithm will be compared with others to prove that it is better than all existing algos.

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